



Foster Hartmann Lafortune





Kavraki



Kress-Gazit





Madhusudan









Expeditions in Computer Augmented Program Engineering

http://excape.cis.upenn.edu/

NSF Expeditions PI Meeting, May 2013



Cornell, Maryland, Michigan, MIT, Penn, Rice, UC Berkeley, UCLA, UIUC



Zdancewic



Vardi



Tripakis



Tabuada Solar-Lezama Seshia



Sangiovanni



Martin

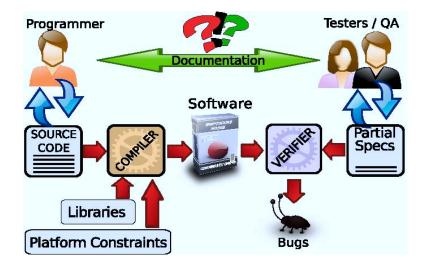


Pappas

Software Design Methodology

- □ What has changed:
 - Programming languages
 - Libraries
 - Verification technology
- □ What has not changed:
 - Programming is done by experts
 - Fully specified by conventional programming
 - Verification phase is distinct from design

Can we leverage modern analysis tools and increased computing power to revolutionize the task of programming?





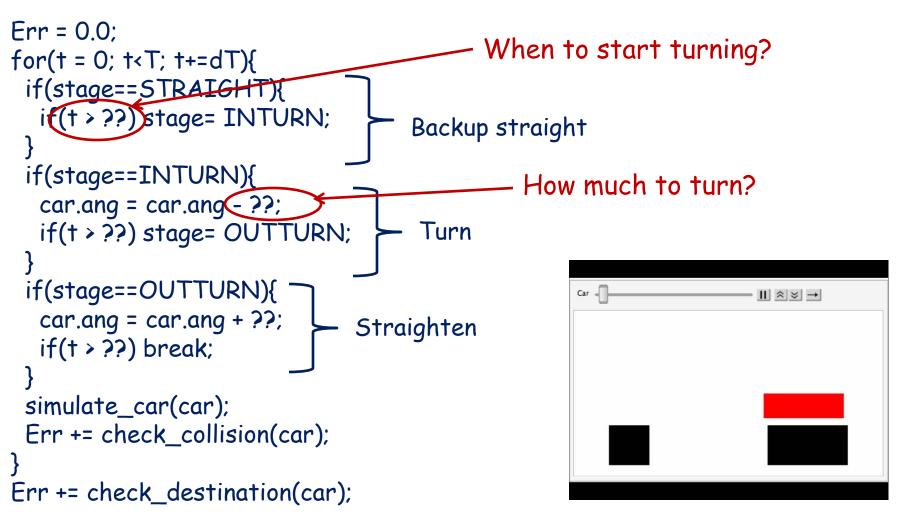
Synthesis: A Plausible Solution?

- □ Classical: Mapping a high-level (e.g. logical) specification to an executable implementation
 - Theoretical foundations: Church (1960s)
 - Derivation of programs from constructive proofs (e.g. Kestrel)
 - Synthesis from temporal logic specs: Clarke/Emerson (1980s)
 - Refinement in model-based design
 - Ongoing progress, but many challenges remain...

Recent shift in focus: Integrating different styles of specifications in a consistent executable

Sketch: Program completion

Ref: Solar-Lezama et al (PLDI 2010)

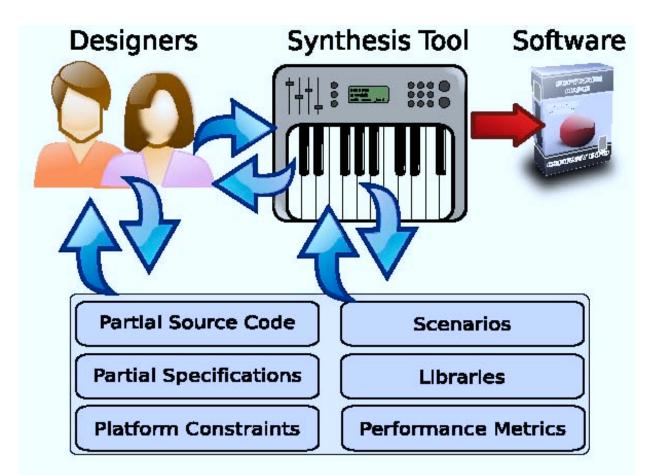


Enables programmers to focus on high-level solution strategy



ExCAPE Vision

Harnessing computation to transform programming: Programming made easier, faster, cheaper



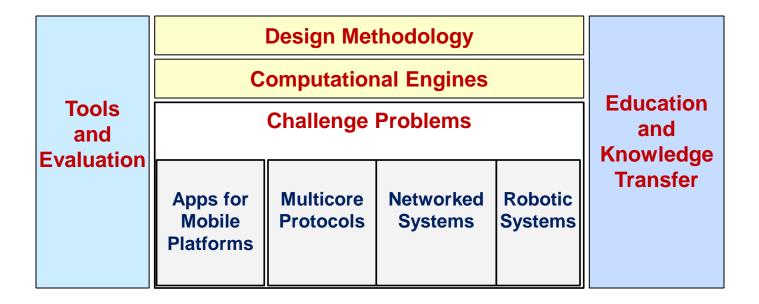
Synthesis Tool: Intelligent Assistance



- Designer expresses "what", possibly using multiple input formats
- Synthesizer discovers new artifacts via integration and completion
- Synthesizer solves computationally demanding problems using advanced analysis tools
- □ Interactive iterative design
- Integrated formal verification



Research Organization



Talk Outline

- Design tool for distributed protocols
- □ Synthesis for programming robots
- □ Synthesis to support online education
- □ Summary of ExCAPE activities

Goal: Simplify Protocol Design

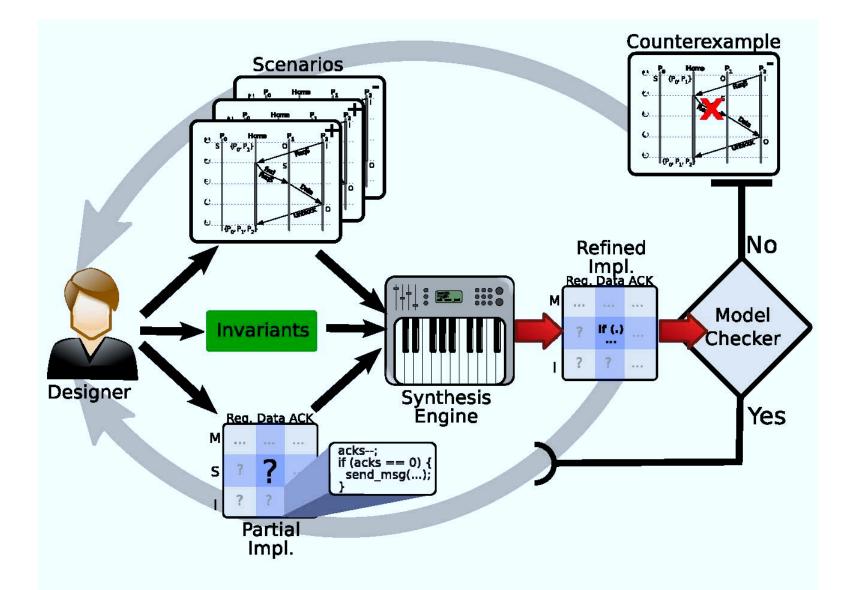
Design challenging due to asynchronous model of communication

- □ Cache coherence protocols, Distributed coordination algorithms
- □ Successful application domain for formal verification / model checking
- Correctness involves both safety and liveness properties
- Proposed solution: Allow programmers flexibility

Skeleton based on Extended-Finite-State-Machines + High-level requirements + Example behaviors



TRANSIT for Distributed Protocol Design



Computational Problem

□ Inputs:

- Variable types and corresponding expression grammar
- For each process,
 - 1. Control states of EFSM
 - 2. List of all variables , input/output messages
 - 3. Set of concrete examples + symbolic constraints
- High-level requirements (invariants and temporal logic formulas)

□ Solution strategy:

- Expression Inference: For each EFSM transition, generate expressions for guards and updates. Solution uses Counter-exampleguided-inductive synthesis using SMT solver Z3
- Check if resulting protocol meets all requirements, using a model checker (Murphi) and if not, report a counter-example

Challenging Case Study: SGI Origin protocol

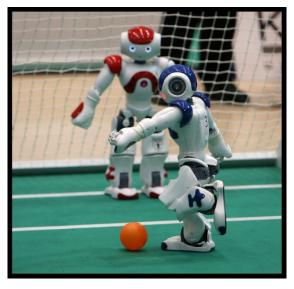
- Source: Laudon and Lenoski; The SGI Origin: A CCNUMA highly scalable server; ISCA 1997
- Directory-based MESI protocol that handles multiple concurrent requests to same requests over unordered network
- Textual description directly leads to protocol skeleton, and symbolic (incomplete) descriptions of most of the transitions
- During debugging, programmer focuses on local fixes of counteexamples and adds concrete examples
- Final iteration required 30 min synthesis time (with 5 Million states explored by Murphi)

 \Box SMT solver / model checker in the loop is feasible for programming 12

Synthesis for Robot Programming

Goal: Allow end-users to program robotic behaviors

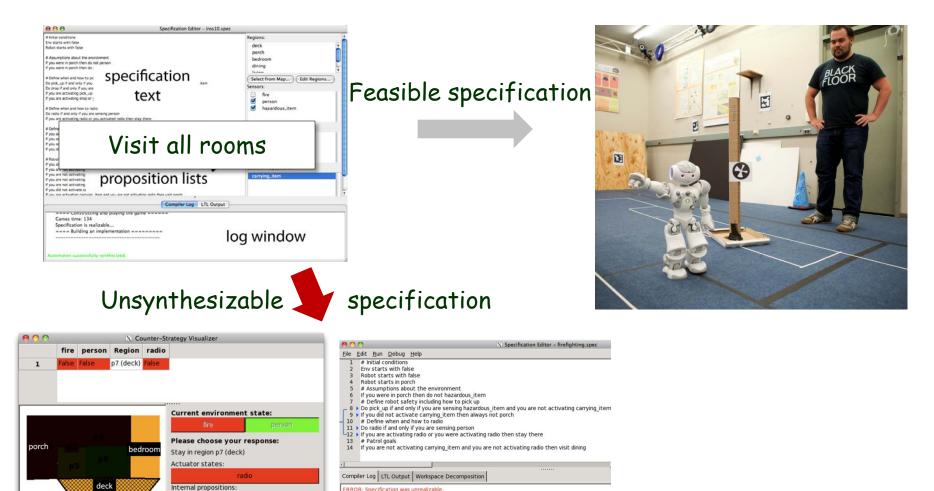




Robotic controllers: Research Challenges

- How to consistently integrate physical constraints, sample trajectories, safety rules, and language/temporal-logic requirements?
- □ How to explain infeasible requirements? How to suggest potential fixes?
- How to program a synthesis engine with completion strategies that take into account the physical and continuous nature of robotics (power, safety, environment traversability)?
- □ How to address optimality and performance?
- □ How to evaluate human-robot interaction?
- How to generate control that ports across different robots (different dynamics, control capabilities, safety considerations)?

LTLMoP: Robot control from structured English



ERROR: Specification was unrealizable.

System is unrealizable because the environment can force a safety violation.

RESULT

Execute Move >>

Currently in step #2

No automator

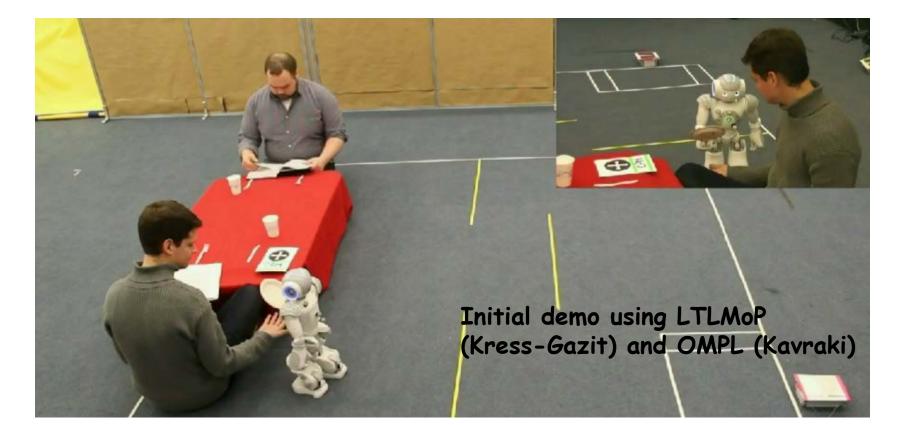
15

Research Results

- Improving the scalability of core engine for mapping Temporal Logic formulas to Controllers: Synthesis with identifiers (Kress-Gazit, Seshia)
- □ Synthesis of cost-optimal plans (Kress-Gazit)
- Motion planning in partially unknown environments (Kavraki, Vardi)
- Synthesis of controllers with robust performance in presence of uncertainties
 Theory of robustness for hybrid systems (Tabuada)
- Accuracy in mapping discrete actions to continuous-time trajectories with durations (Kress-Gazit)
- □ Automatic generation of environment assumptions (Alur, Topcu)

Ongoing Case Study: Robotic Waiter

- Challenges: Scalability (items, costumers), uncertainty in sensing and actuation, optimality of behavior, fault recovery
- Future plans: exploit symmetries, robust synthesis, task specific abstractions



Synthesis for Online Education

Emerging opportunity: MOOCs



coursera

□ Challenge: Personalized feedback on assignments

- Manual feedback by TAs (not scalable)
- Grading by peers (not reliable)
- Evaluation on test cases (how to translate failed tests to errors?)

□ Application for ExCAPE tools for synthesis

- Introductory programming assignments (Solar-Lezama, MIT)
- Scheduling problems in Embedded Systems course (Seshia, UC Berkeley)
- DFA construction in Theory of Computation (with Hartmann, UC Berkeley) see automatatutor.com

Sample Problem: Derivative of a Polynomial

```
def computeDeriv(poly):
    result = []
    for i in range(len(poly)):
        result += [i * poly[i]]
    if len(poly) == 1:
        return result # return [0]
    else:
        return result[1:] # remove the leading 0
```

Autograder Output on a Student Solution

```
def computeDeriv(poly):
 1
 2
3
4
5
6
        deriv = []
        zero = 0
        if (len(poly) == 1):
            return deriv
        for e in range(0,len(poly)):
 7
            if(poly[e]==0):
 8
                 zero += 1
 9
            else:
10
                 deriv.append(poly[e]*e)
11
```

return deriv

12

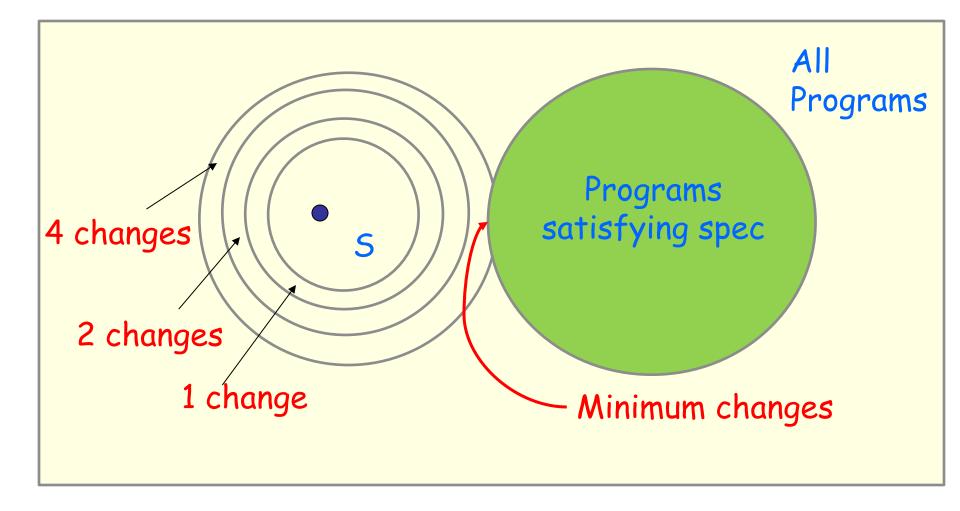
Student Solution + Reference Solution + Error Model



The program requires 3 changes:

- In the return statement **return deriv** in **line 5**, replace **deriv** by **[0**].
- In the comparison expression (poly[e] == 0) in line 7, change (poly[e] == 0) to False.
- In the expression range(0, len(poly)) in line 6, replace 0 by 1.

Computational Engine: Sketch Synthesis Tool



S: Student Solution

Autograder Experiments (MIT 6.00)

	Generated		
TestSet	Feedback	Percentage	AvgTime(s)
268	218	81.34%	2.49
344	185	53.78%	2.65
103	88	85.44%	12.95
13	6	46.15%	3.35
52	17	32.69%	29.57
918	753	82.03%	12.42
541	167	30.87%	4.78
1756	860	48.97%	4.14
2875	1693	58.89%	3.58
2938	2271	77.30%	10.59
2988	2052	68.67%	17.13
351	171	48.72%	9.08
218	98	44.95%	22.09
	268 344 103 13 52 918 541 1756 2875 2938 2938 2988 351	TestSetFeedback2682183441851038813652179187535411671756860287516932938227129882052351171	TestSetFeedbackPercentage26821881.34%34418553.78%1038885.44%13646.15%521732.69%91875382.03%54116730.87%175686048.97%2875169358.89%2938227177.30%2988205268.67%35117148.72%

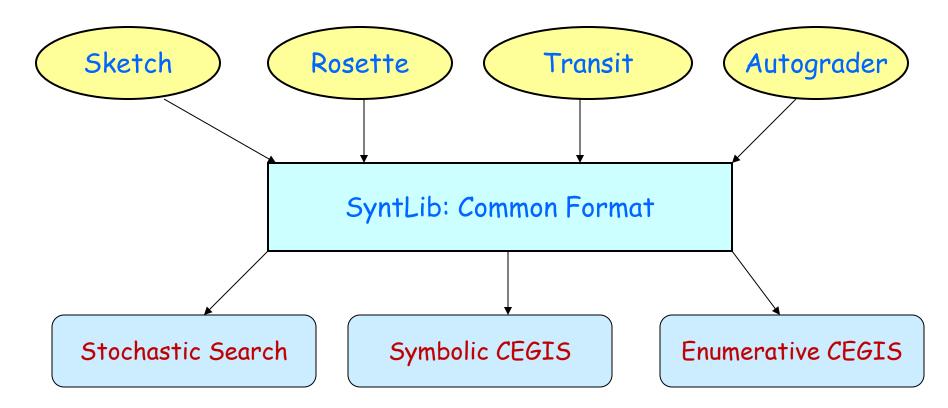
64%

9.9s

Theory / Methodology / Tools / Application domains

- □ Rosette: Framework for developing synthesis-enhanced DSLs (Bodik)
- Enhancements to Sketch to support modularity (Solar-Lezama)
- Bridging the gap between reactive synthesis and supervisory control (Lafortune, Tripakis, Vardi)
- Verified LLVM Infrastructure (Martin, Zdancewic)
- □ Platform-based design for software synthesis (Sangiovanni-Vincetelli)
- □ Synthesis of logic for avoiding concurrency bugs (Lafortune)
- □ Component-based synthesis for probabilistic systems (Vardi)
- □ Theory of regular functions for quantitative analysis (Alur)
- □ Automated cloud configuration using synthesis(Alur, Loo, Parthasarathy)
- □ Route Shepherd for configuration of routing protocols (Loo)
- □ Synthesis of control + scheduling for wireless control networks (Pappas)
- Programming for mobile platforms (Foster, Solar-Lezama)
- □ User studies for improving programming notations (Hartmann)

Template-based Synthesis Modulo Theories



Based on input format for SMTLib 2
Problem: Given a formula φ in an SMT theory with an extra function symbol f, and context-free language L for templates, find an expression e in L such that φ[f/e] is valid
Basis for synthesis competition (to be held at CAV 2014)

Education and Outreach

 ExCAPE Summer School: June 13—16, Berkeley; 125 registrants Tutorials: Reactive synthesis (Vardi) Constraint-based program synthesis (Bodik) Synthesis for cyber-physical systems (Tabuada) + Talks

□ ExCAPE Webinar: Monthly talks on diverse topics

Sponsored workshops SYNT (at CAV 2013, by Solar-Lezama) Synthesis for robotics (at RSS 2013, by Kavraki and Kress-Gazit) Special session on synthesis at ACC 2013 (by Lafortune)

Graduate course at Berkeley: Program synthesis for everyone

□ K-12 programs: CURIE @ Cornell

Rotating Postdoc Program

□ Each ExCAPE postdoc has two mentors, at two different institutions

Year 2012-13: Ruediger Ehlers (Robotics) Mentors: Kress-Gazit (Cornell), Seshia (UC Berkeley) □ For the upcoming year: Xiaokang Qiu (UIUC), Apps for mobile platforms Mentors: Foster (Maryland), Solar-Lezama (MIT) Indranil Saha (UCLA), Robotics Mentors: Pappas (Penn), Seshia (UC Berkeley) Christos Stergiou (UC Berkeley), Multicore protocols Mentors: Martin (Penn), Tripakis (UC Berkeley) Damian Zufferey (IST Austria), Networked systems

Mentors: Loo (Penn), Parthasarathy (UIUC)





Paradigm shift in synthesis:

Old: Allow more concise, high-level description New: Designer uses multiple, natural formats,

Synthesis tool assists in discovering tricky logic

Paradigm shift in design tools:
 Old : Any compiler transformation must be polynomial-time
 New: Computational intractability not a show-stopper

 Common theme: Guided search in a space of programs to find one that meets multiple design goals
 A bit like model checking, but can be interactive!